

RUTGERS



1-Sperner hypergraphs and new characterizations of threshold graphs

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Joint work with

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Background and motivation

Hypergraphs

Hypergraph: a pair $\mathcal{H} = (V, \mathcal{E})$ where

- ▶ V is a finite set of **vertices**
- ▶ \mathcal{E} is a set of subsets of V , called **hyperedges**

Example:

- ▶ $V = \{1, 2, 3, 4\}$
- ▶ $\mathcal{E} = \{\{1, 2\}, \{1, 3\}, \{1, 4\}, \{1, 2, 3\}, \{2, 3, 4\}\}$

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$\{1, 2, 3\}$

$\{2, 3, 4\}$

$\{1, 2\}$

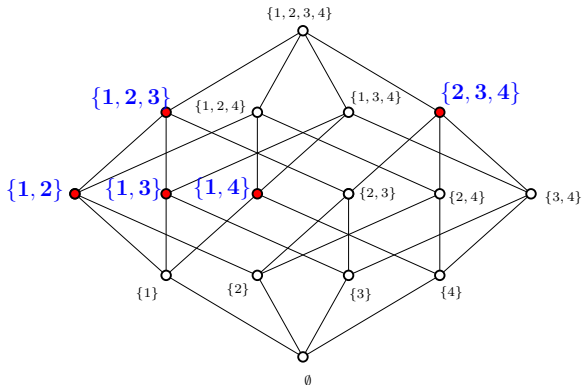
$\{1, 3\}$

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Hypergraphs

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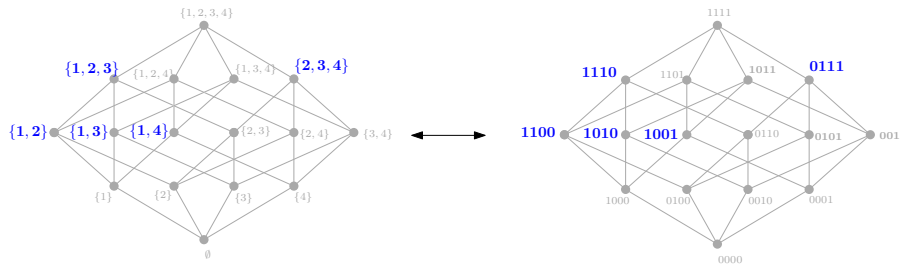
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set $S \subseteq V$ \longleftrightarrow $\chi^S \in \{0, 1\}^V$
characteristic vector of S

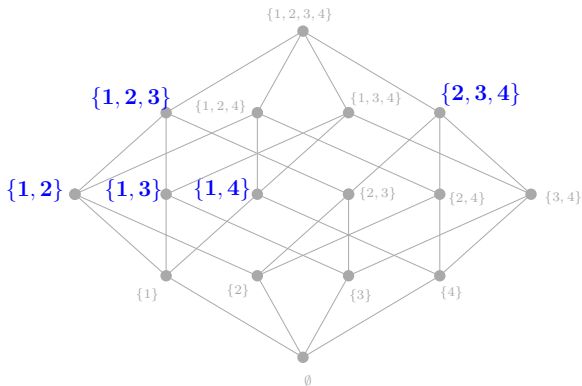
Hypergraphs

An **independent set** in a hypergraph $\mathcal{H} = (V, \mathcal{E})$ is a set $X \subseteq V$ containing no hyperedge of \mathcal{H} .

A set $X \subseteq V$ is **dependent** if it is not independent.

Equivalently: if it contains a hyperedge.

Example:

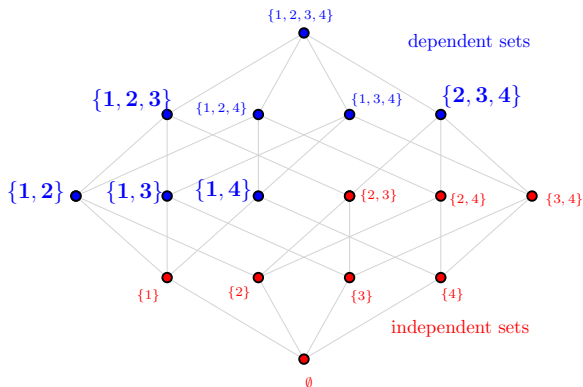


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Example:



Threshold hypergraphs

A hypergraph \mathcal{H} is **threshold** if $\exists w : V \rightarrow \mathbb{Z}_{\geq 0}, \exists t \in \mathbb{Z}_{\geq 0}$ such that

for all $X \subseteq V$:

$$X \text{ is a dependent set of } \mathcal{H} \iff w(X) \geq t$$

$$w(X) := \sum_{x \in X} w(x).$$

- ▶ **Geometric interpretation:** *there is a hyperplane separating the characteristic vectors of independent sets from the characteristic vectors of dependent sets.*

Threshold hypergraphs - example

$\exists w : V \rightarrow \mathbb{Z}_{\geq 0}, \exists t \in \mathbb{Z}_{\geq 0}$ such that
for all $X \subseteq V$: X is a dependent set of $\mathcal{H} \iff w(X) \geq t$.

Example:

$\{1, 2, 3\}$

$\{2, 3, 4\}$

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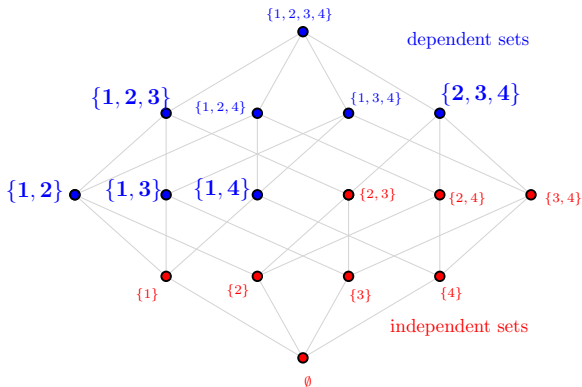
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Threshold hypergraphs - example

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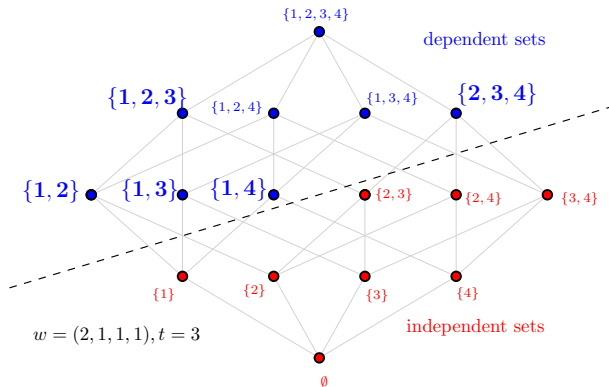


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Example:



Threshold hypergraphs

Some historical remarks:

- ▶ Threshold hypergraphs were defined in the uniform case by **Golumbic** in 1980 and studied further by **Reiterman**, **Rödl**, **Šiňajová**, **Tůma** in 1985.
- ▶ In their full generality, the concept of threshold hypergraphs is equivalent to that of **threshold monotone Boolean functions** studied, e.g., by **Muroga** in 1971.

Threshold hypergraphs

- ▶ A polynomial time recognition algorithm for threshold monotone Boolean functions represented by their complete DNF was given by [Peled and Simeone](#) in 1985.
- ▶ The algorithm is based on linear programming and implies the existence of a **polynomial time recognition algorithm for threshold hypergraphs**.

Threshold hypergraphs arising from graphs

Several classes of threshold hypergraphs arising from graphs were studied in the literature:

G – a graph

hyperedges	studied by	resulting graph class
edges (of G)	Chvátal-Hammer, 1977	threshold graphs
vertex covers	–	threshold graphs
dominating sets	Benzaken-Hammer, 1978	domishold graphs
total dom. sets	Chiarelli-M., 2014-15	total dom. graphs
connected dom. sets	Chiarelli-M., 2014	connected dom. graphs
maximal cliques	???	???

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maximal cliques	this talk	threshold graphs

Sperner hypergraphs

A hypergraph is said to be **Sperner** (or: a **clutter**) if no hyperedge contains another one, that is,

if $e, f \in \mathcal{E}$ and $e \subseteq f$ implies $e = f$.

Example:

▶ $V = \{1, 2, 3, 4\}$

▶ $\mathcal{E} =$

$$\{1, 2, 3\}$$

$$\{2, 3, 4\}$$

$$\{1, 2\}$$

$$\{1, 3\}$$

$$\{1, 4\}$$

not Sperner since $\{1, 2\} \subset \{1, 2, 3\}$

Dually Sperner hypergraphs

Sperner hypergraphs can be equivalently defined as the hypergraphs such that

$$e \neq f \Rightarrow \min\{|e \setminus f|, |f \setminus e|\} \geq 1.$$

This point of view motivated Chiarelli and M. to define in 2014 a hypergraph \mathcal{H} to be **dually Sperner** if

$$e \neq f \Rightarrow \min\{|e \setminus f|, |f \setminus e|\} \leq 1.$$

Dually Sperner hypergraphs

$$e \neq f \Rightarrow \min\{|e \setminus f|, |f \setminus e|\} \leq 1.$$

Example:

The hypergraph from the previous example is dually Sperner:

► $\mathcal{E} =$

$\{1, 2, 3\}$

$\{2, 3, 4\}$

$\{1, 2\}$

$\{1, 3\}$

$\{1, 4\}$

The following hypergraph is not dually Sperner:

► $V = \{1, 2, 3, 4\}$

► $\mathcal{E} = \{\{1, 2\}, \{3, 4\}\}$

Dually Sperner hypergraphs

Theorem (Chiarelli-M. 2014)

Every dually Sperner hypergraph is threshold.

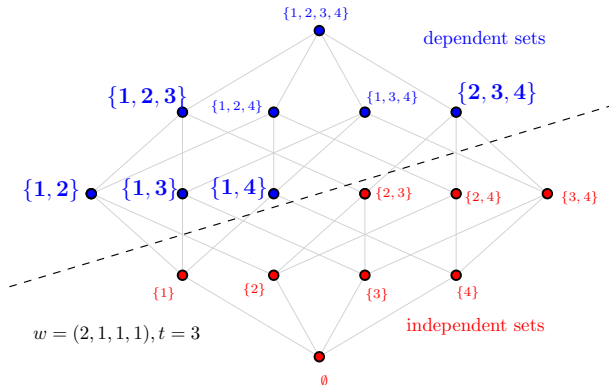
Chiarelli and M. applied dually Sperner hypergraphs to characterize two classes of graphs related to separation of total, resp. connected dominating sets.

Threshold hypergraphs

A hypergraph \mathcal{H} is **threshold** if

$\exists w : V \rightarrow \mathbb{Z}_{\geq 0}, \exists t \in \mathbb{Z}_{\geq 0}$ such that
for all $X \subseteq V$:

X is an dependent set of $\mathcal{H} \iff w(X) \geq t.$



Threshold hypergraphs

It follows from the definition of threshold hypergraphs that **only minimal hyperedges matter** for the thresholdness property of a given hypergraph.

Example:

$\{1, 2, 3\}$

$\{2, 3, 4\}$

$\{1, 2\}$

$\{1, 3\}$

$\{1, 4\}$

is threshold if and only if

$\{2, 3, 4\}$

$\{1, 2\}$

$\{1, 3\}$

$\{1, 4\}$

is threshold.

1-Sperner hypergraphs

Since dually Sperner hypergraphs are threshold, we focus on the family of hypergraphs that are **both Sperner and dually Sperner**.

We call such hypergraphs **1-Sperner**.

A hypergraph \mathcal{H} is 1-Sperner if and only if

$$e \neq f \Rightarrow \min\{|e \setminus f|, |f \setminus e|\} = 1.$$

1-Sperner hypergraphs

$$e \neq f \Rightarrow \min\{|e \setminus f|, |f \setminus e|\} = 1.$$

Example:

The hypergraph from the previous example is not 1-Sperner, since it is not Sperner:

► $\mathcal{E} =$

$\{1, 2, 3\}$

$\{2, 3, 4\}$

$\{1, 2\}$

$\{1, 3\}$

$\{1, 4\}$

Deleting the hyperedge $\{1, 2, 3\}$ results in a 1-Sperner hypergraph:

► $V = \{1, 2, 3, 4\}$

► $\mathcal{E} = \{\{1, 2\}, \{1, 3\}, \{1, 4\}, \{2, 3, 4\}\}$

Our results

Our results

1. A composition result for 1-Sperner hypergraphs.
2. Its consequences.
3. New characterizations of threshold graphs.

An operation preserving 1-Spernerness

Gluing of hypergraphs

$\mathcal{H}_1 = (V_1, \mathcal{E}_1)$ and $\mathcal{H}_2 = (V_2, \mathcal{E}_2)$ – vertex-disjoint hypergraphs

z – a new vertex

The **gluing** of \mathcal{H}_1 and \mathcal{H}_2 is the hypergraph

$$\mathcal{H} = \mathcal{H}_1 \odot \mathcal{H}_2$$

such that

$$V(\mathcal{H}) = V_1 \cup V_2 \cup \{z\}$$

and

$$E(\mathcal{H}) = \{\{z\} \cup e \mid e \in \mathcal{E}_1\} \cup \{V_1 \cup e \mid e \in \mathcal{E}_2\}.$$

Incidence matrices

The operation of gluing can be visualized easily in terms of **incidence matrices**.

Every hypergraph $\mathcal{H} = (V, \mathcal{E})$ with $V = \{v_1, \dots, v_n\}$ and $\mathcal{E} = \{e_1, \dots, e_m\}$

can be represented with its **incidence matrix** $A^{\mathcal{H}} \in \{0, 1\}^{m \times n}$:
rows are indexed by hyperedges of \mathcal{H} ,
columns are indexed by vertices of \mathcal{H} ,

and

$$A_{i,j}^{\mathcal{H}} = \begin{cases} 1, & \text{if } v_j \in e_i; \\ 0, & \text{otherwise.} \end{cases}$$

Gluing of hypergraphs

If $\mathcal{H} = \mathcal{H}_1 \odot \mathcal{H}_2$ then

$$A^{\mathcal{H}_1 \odot \mathcal{H}_2} = \begin{pmatrix} \mathbf{1}^{m_1,1} & A^{\mathcal{H}_1} & \mathbf{0}^{m_1,n_2} \\ \mathbf{0}^{m_2,1} & \mathbf{1}^{m_2,n_1} & A^{\mathcal{H}_2} \end{pmatrix}.$$

Example:

$$\begin{matrix} A^{\mathcal{H}_1} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \\ A^{\mathcal{H}_2} = \begin{pmatrix} 1 & 1 & 0 \\ 1 & 0 & 1 \\ 0 & 1 & 1 \end{pmatrix} \end{matrix} \longrightarrow A^{\mathcal{H}_1 \odot \mathcal{H}_2} = \begin{pmatrix} z & & & & & \\ \hline 1 & 1 & 0 & | & 0 & 0 & 0 \\ 1 & 0 & 1 & | & 0 & 0 & 0 \\ \hline 0 & 1 & 1 & | & 1 & 1 & 0 \\ 0 & 1 & 1 & | & 1 & 0 & 1 \\ 0 & 1 & 1 & | & 0 & 1 & 1 \end{pmatrix}$$

Gluing of hypergraphs

Proposition

For every pair $\mathcal{H}_1 = (V_1, \mathcal{E}_1)$ and $\mathcal{H}_2 = (V_2, \mathcal{E}_2)$ of vertex-disjoint 1-Sperner hypergraphs,

their gluing $\mathcal{H}_1 \odot \mathcal{H}_2$ is a 1-Sperner hypergraph,

unless $\mathcal{E}_1 = \{V_1\}$ and $\mathcal{E}_2 = \{\emptyset\}$

(in which case $\mathcal{H}_1 \odot \mathcal{H}_2$ is not Sperner).

$$\begin{array}{l} A^{\mathcal{H}_1} = (1 \quad 1) \\ A^{\mathcal{H}_2} = (0 \quad 0 \quad 0) \end{array} \longrightarrow A^{\mathcal{H}_1 \odot \mathcal{H}_2} = \left(\begin{array}{c|cc|ccc} z & 1 & 1 & 0 & 0 & 0 \\ \hline 0 & 1 & 1 & 0 & 0 & 0 \end{array} \right)$$

A composition result for 1-Sperner hypergraphs

We show that every nontrivial 1-Sperner hypergraph can be generated this way.

We say that a gluing of two vertex-disjoint hypergraphs $\mathcal{H}_1 = (V_1, \mathcal{E}_1)$ and $\mathcal{H}_2 = (V_2, \mathcal{E}_2)$ is **safe** unless $\mathcal{E}_1 = \{V_1\}$ and $\mathcal{E}_2 = \{\emptyset\}$.

Theorem

A hypergraph \mathcal{H} is 1-Sperner if and only if

*it either **has no vertices** (that is, $\mathcal{H} \in \{(\emptyset, \emptyset), (\emptyset, \{\emptyset\})\}$)*

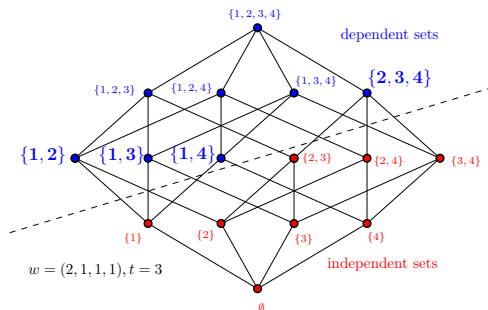
or it is a safe gluing of two smaller 1-Sperner hypergraphs.

Consequences of the structural result

Consequences

Using the composition result for 1-Sperner hypergraph, we obtain the following:

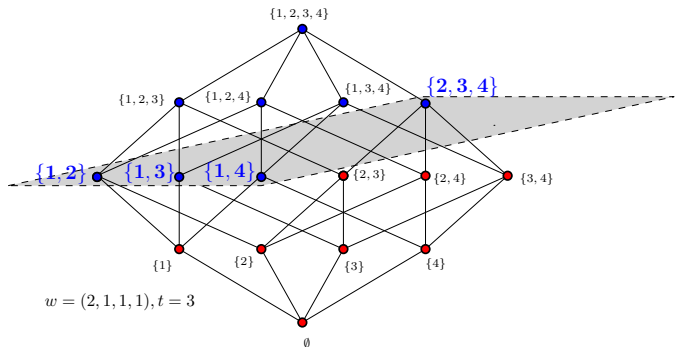
1. An alternative proof of the fact that every 1-Sperner hypergraph is threshold.



- ▶ Unlike the previous proof establishing thresholdness of dually Sperner hypergraphs (due to Chiarelli-M.), this proof is constructive and builds a separating hyperplane of a given 1-Sperner hypergraph.

Consequences

2. A proof of the fact that every 1-Sperner hypergraph is equilizable.



- ▶ Equilizable hypergraphs form a generalization of **equistable graphs** (introduced in 1980 by Payan and studied afterwards in over 10 papers).

Consequences

3. An upper bound on the size of 1-Sperner hypergraphs:

Proposition

For every 1-Sperner hypergraph $\mathcal{H} = (V, \mathcal{E})$ with $\mathcal{E} \neq \{\emptyset\}$, we have $|\mathcal{E}| \leq |V|$.

- ▶ Proof idea: the characteristic vectors of the hyperedges are linearly independent in \mathbb{R}^V .

Can we prove a similar **lower bound**?

Consequences

- ▶ **universal** vertex: a vertex contained in all hyperedges
- ▶ **isolated** vertex: a vertex contained in no hyperedges
- ▶ two vertices u, v are **twins** if they are contained in exactly the same hyperedges

Adding universal vertices, isolated vertices, or twin vertices preserves the 1-Sperner property, while

- ▶ keeping the number of hyperedges unchanged and
- ▶ increasing the number of vertices.

Consequently, there is no lower bound on the number of hyperedges of a 1-Sperner hypergraph in terms of the number of vertices.

However . . .

Consequences

4. A lower bound on the size of 1-Sperner hypergraphs without universal, isolated, and twin vertices:

Proposition

For every 1-Sperner hypergraph $\mathcal{H} = (V, \mathcal{E})$ with $|V| \geq 2$ and without universal, isolated, and twin vertices, we have

$$|\mathcal{E}| \geq \left\lceil \frac{|V| + 2}{2} \right\rceil.$$

This bound is sharp.

New characterizations of threshold graphs

Threshold graphs

A **threshold graph** is a threshold hypergraph in which all hyperedges are of size 2.

- ▶ Threshold graphs were introduced by **Chvátal and Hammer** in the 1970s and were studied in numerous papers (and in a monograph by **Mahadev and Peled** from 1995).
- ▶ Threshold graphs have many different characterizations.

Threshold graphs

Theorem (Chvátal and Hammer, 1977)

A graph G is threshold if and only if it is $\{P_4, C_4, 2K_2\}$ -free.

Theorem (Chvátal and Hammer, 1977)

A graph G is threshold if and only if $V(G) = K \cup I$

where K is a clique, I is an independent set, $K \cap I = \emptyset$, and

*there exists an ordering v_1, \dots, v_k of I such that $N(v_i) \subseteq N(v_j)$
for all $1 \leq i < j \leq k$.*

Clique hypergraphs of graphs

Given a graph G , the **clique hypergraph of G** is the hypergraph $\mathcal{C}(G)$ with vertex set $V(G)$ in which the hyperedges are exactly **the maximal cliques of G** .

Theorem (Berge, 1989)

*The clique hypergraphs of graphs are exactly those Sperner hypergraphs \mathcal{H} that are also **normal** (or: **conformal**), that is,*

for every set $X \subseteq V(\mathcal{H})$ such that every pair of elements in X is contained in a hyperedge, there exists a hyperedge containing X .

A necessary condition for thresholdness

A hypergraph is **k -summable** if it has k (not necessarily distinct) independent sets A_1, \dots, A_k and k (not necessarily distinct) dependent sets B_1, \dots, B_k such that

$$\sum_{i=1}^k \chi^{A_i} = \sum_{i=1}^k \chi^{B_i}.$$

If a graph is k -summable for some $k \geq 2$, then it cannot be threshold.

A hypergraph is **k -asummable** if it is not k -summable.

A necessary condition for thresholdness

Theorem

A hypergraph is threshold if and only if it is k -asummable for all k .

- ▶ A restatement of the analogous characterization of threshold Boolean functions proved in 1961 independently by Chow and Elgot.

Corollary

Every threshold hypergraph is 2-asummable.

1-Sperner, threshold, 2-asummable

In general:

$$1\text{-Sperner} \Rightarrow \text{threshold} \Rightarrow 2\text{-asummable}$$

and none of the implications can be reversed.

In the class of **conformal** Sperner hypergraphs, all these three notions coincide.

Moreover, they **exactly characterize threshold graphs**.

New characterizations of threshold graphs

Theorem

For every graph G , the following statements are equivalent:

- (1) G is threshold.
- (2) The clique hypergraph $\mathcal{C}(G)$ is 1-Sperner.
- (3) The clique hypergraph $\mathcal{C}(G)$ is threshold.
- (4) The clique hypergraph $\mathcal{C}(G)$ is 2-asummable.

clique hypergraph $\mathcal{C}(G)$

\rightsquigarrow **independent set hypergraph $\mathcal{I}(G)$**

in (2), (3), (4) also ok

(since the class of threshold graphs is closed under taking complements)

Summary

- ▶ We introduced a new class of hypergraphs, the class of **1-Sperner hypergraphs**:
 $e \neq f \Rightarrow \min\{|e \setminus f|, |f \setminus e|\} = 1.$
- ▶ We proved a **structural theorem** for 1-Sperner hypergraphs and examined several of its **consequences**, including **bounds** on the size of 1-Sperner hypergraphs and a **new, constructive proof** of the fact that every 1-Sperner hypergraph is threshold.
- ▶ Within the class of **normal Sperner hypergraphs**:
1-Sperner \Leftrightarrow threshold \Leftrightarrow 2-asummable
- ▶ New characterizations of the class of **threshold graphs**.

Thank you!

Hvala!